

Now, let us generalize Mass Distribution Principle.

For $x \in \mathbb{R}^d$, let $Q_n(x)$ be b^{-n} b -adic cube containing x , as before.

Lemma (Billingsley)

Let K be a Borel set in \mathbb{R}^d , μ -finite Borel measure on \mathbb{R}^d and $\mu(K) > 0$. If for some $\beta > 2 \geq 0$

$$\underline{d} \leq \lim_{n \rightarrow \infty} \frac{\log \mu(Q_n(x))}{-n \log b} \leq \beta \quad \forall x \in K, \quad \text{then} \quad \underline{d} \leq \text{Hdim } K \leq \beta.$$

Pf. Take $\underline{d}_1 < \underline{d}$. By the left inequality,

$\lim_{n \rightarrow \infty} b^{2n} \mu(Q_n(x)) = 0 \quad \forall x \in K$. By singularity lemma applied to $h(x) = x^{2^i}$, $m_{\beta_1}(K) = \infty$, so $\text{Hdim } K \geq \underline{d}_1, \forall \underline{d}_1 < \underline{d}$.

In the other direction, for $\beta_1 > \beta$, we have

$\lim_{n \rightarrow \infty} b^{\beta_1 n} \mu(Q_n(x)) = \infty$ for all $x \in K$.

Fix $\varepsilon > 0$, and for any $x \in K$ let $n(x)$ be the smallest n such that $b^{\beta_1 n} \mu(Q_n(x)) > \varepsilon$ and $\forall d \leq b^{-n} \leq \varepsilon$. $\{Q_{n(x)}(x) : x \in K\}$ is a cover of K . select a non-intersecting subcover $\{Q_k\}$. Then $\text{diam } Q_k \leq \varepsilon$ (by the choice of $n(x)$),

$\sum (\text{diam } Q_k)^{\beta_1} \leq \sum \mu(Q_k) \leq \mu(\mathbb{R}^d)$. Thus $m_{\beta_1, \varepsilon}(K) \leq \mu(\mathbb{R}^d)$. Take $\varepsilon \rightarrow 0$, to get that $m_{\beta_1}(K) < \infty$.

An application: dimension of lacunary set.

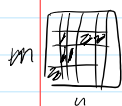
Let us show that $\text{Hdim } A_S = \underline{d}(S)$.

Let us define μ to be - measure giving equal weight to $2^{-n} (S \cap (1, \dots, n))$ to each of the n -th generation intervals of C_n . For $x \in C$,

$\lim_{n \rightarrow \infty} \frac{\log \mu(Q_n(x))}{\log 2^{-n}} = \underline{d}(S)$. So, by Billingsley, $\text{Hdim } A_S = \underline{d}(S)$.

Another application: self-similar sets

Let D be an $n \times m$ pattern on the square:



Formally, $D \subset \{0, \dots, n-1\} \times \{0, \dots, m-1\}$

Repeat the pattern on every non-removed rectangle.

$K(D) := \{ \sum_{k=0}^{\infty} (a_k n^{-k}, b_k m^{-k}) : (a_k, b_k) \in D \}$.

Most famous: McMullen set, corresponding

to $D = \{(1,0), (1,1), (2,0)\}$ in 3×2 world.

Lemma (Easy). Assume $n > m$ and that every row contains a chosen rectangle (i.e. $\forall j \exists i: (i,j) \in D$).

Then $\text{Mdim } K(D) = 1 + \log_n \#(D)$.

Exercise: What happens when not every row contains a chosen rectangle?

Proof. Let $v := \#(D)$ and let us look at the j -th stage of the

Proof. Let $r := \#(D)$ and let us look at the j -th stage of the construction. It consists of r elongated rectangles of width n^{-j} and height m^{-j} . Let $k = \lfloor \frac{\log n}{\log m} j \rfloor$. Then we need



$\frac{m^{-j}}{n^{-j}} = m^{k-j}$ squares of size m^{-k} to cover each such rectangle, and we cannot use less than m^{k-j} (because every row has a rectangle in it, thus the height of $k(D) \cap$ any given rectangle is still m^{-j}).

$$N(k(D), m^{-k}) \approx r j m^{k-j},$$

$$M \dim k(D) = \lim_{k \rightarrow \infty} \frac{\log r + \log n^{k-j} + \log m^{k-j}}{k \log m} =$$

$$\frac{\log r}{\log n} + \frac{\log n}{\log n} - \frac{\log m}{\log n} = 1 + \log_n \frac{r}{m} \in$$

So, for McMullen set,
 $\text{Hdim } k(D) = 1 + \log_{\frac{2}{3}} \frac{r}{m}$.

But

Thm. In the assumptions of the lemma,
 $\text{Hdim}(k(D)) = \log_m(\sum r(j) \log_m m)$ where

$r(j) = \# \{i : (i, j) \in D\}$ - number of rectangles lying in a given row.

It is easy to see that $\text{Hdim}(k(D)) \leq M \dim(k(D))$, with equality reached iff $r(j)$ does not depend on j .

For McMullen set, $\dim k(D) = \log_2(1 + 2^{\log_2 2})$

Pt Instead of b -adic squares, we use more appropriate approximate squares

$$Q_k(x, y) = \left\{ (x', y') : \begin{array}{l} x' = \sum x'_j n^{-j} \quad x'_j = x_j, j \leq k \\ y' = \sum y'_j m^{-j} \quad y'_j = y_j, j \leq k \end{array} \right\}$$

We expand x in base n | y in base m !

It is almost a square, since $m^{-k} \leq n^{-(k+1)} \leq n m^{-k}$, and $\text{diam } Q_k \subseteq m^{-k}$ (up to a factor of $\sqrt{2}n$).

Given a probability vector \vec{p} with coordinates $p(d), d \in D$, define the probability measure $\mu_{\vec{p}}$ on $k(D)$ by the usual iterative process. Assume also that

$\tilde{p}(d_1) = \tilde{p}(d_2)$ if d_1, d_2 are in the same row (i.e. $y_{d_1} = y_{d_2}$)

Thus if $d = (x_j)_{j=1}^{\infty}$, $x_j \in \{0, \dots, n-1\}$, $j \in \{0, \dots, m-1\}$, $\tilde{p}(d) = \tilde{p}(j)$, $\sum r(j) p(j) = 1$.

$$\mu_{\vec{p}} Q_k(x, y) = \prod_{j=1}^k p(y_j) \prod_{j=2^k+1}^{\infty} r(y_j)$$

(since $Q_k(x, y)$ contains exactly $\prod_{j=2^k+1}^{\infty} r(y_j)$ rectangles of level k , each of the same measure).

$$\log(\mu_{\vec{p}}(Q_k(x, y))) = \sum_{j=1}^k \log p(y_j) + \sum_{j=2^k+1}^{\infty} \log r(y_j)$$

Then, since (y_j) are i.i.d. w.r.t $\mu_{\vec{p}}$, by SLLN, (or by ergodic thm, discuss next)

$$\lim_{k \rightarrow \infty} \frac{1}{k} \log(\mu_{\vec{p}}(Q_k(x, y))) = \sum_j r(j) p(j) \log p(j) + (1 - \sum_j r(j) p(j)) \log r(j) \mu_{\vec{p}} \text{ a.e.}$$

and, by a version of Billingsley's Lemma for approximate squares,

$$\text{Hdim } k(D) \geq \frac{1}{\log m} \sum_j r(j) p(j) \left(\log \frac{1}{p(j)} + \log(r(j) \mu_{\vec{p}}) \right)$$

To maximize it pick $\dots 2-1 \frac{m-1}{m-1} \dots 2$

$$Hdim k(D) \geq \frac{1}{\log m} \sum_{j \in D} \left(\log \frac{1}{r(j)} + \log (r(j)^{2^{-1}}) \right).$$

To maximize it pick

$$p(j) = \frac{1}{z} r(j)^{2^{-1}}, \quad z = \sum_{j \in D} r(j)^{2^{-1}} = \sum_{j=0}^{m-1} r(j)^{2^{-1}}.$$

Then $Hdim k(D) \geq \log_m z$, and, thus we obtain the right lower bound for $Hdim k(D)$.
For the upper bound, it

$$S_k(x, y) := \sum_{j=1}^k \log r(y_j).$$

$$\begin{aligned} \frac{1}{k} S_k(x, y) & \text{ is uniformly bounded,} \\ \log \mu(Q_k(x, y)) &= \sum_{j=1}^k \log \frac{1}{z} r(y_j)^{2^{-1}} + \left(\sum_{j=1}^k \log r(y_j) - \sum_{j=1}^{2^k} \log r(y_j) \right) = \\ &= -\sum_{j=1}^k \log z + (k-1) S_k(x, y) + S_k(x, y) - S_{2^k}(x, y), \text{ or} \\ \log \mu(Q_k(x, y)) + k \log z &= 2 S_k(x, y) - S_{2^k}(x, y). \end{aligned}$$

$$\text{so: } \frac{1}{2k} \log \mu(Q_k(x, y)) + \frac{1}{2} \log z = \frac{S_k(x, y)}{k} - \frac{S_{2^k}(x, y)}{2^k}.$$

Now do the telescoping series to get that, since

$$\sum \frac{S_{2^{-j}}(x, y) - S_{2^{-j+1}}(x, y)}{2^{-j}} \text{ converges,}$$

$$\lim \left(\frac{S_k(x, y)}{k} - \frac{S_{2^k}(x, y)}{2^k} \right) \geq 0. \Rightarrow$$

$$\lim \left(\log \mu(Q_k(x, y)) + k \log z \right) \geq 0 \Rightarrow$$

$$\lim \frac{\log \mu(Q_k(x, y))}{-k} \leq \log z. \text{ so, since}$$

$$\text{diam } Q_k(x, y) \asymp m^{-k}, \quad Hdim(k(D)) \leq \log_m z$$

(we use again an approximate square version of Billingsley's Lemma)